

Derived categories and functors

Plan

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- ② Localisation
- ③ Derived functors

§1 Recollections and motivation

- A **triangulated category** \mathcal{T} is an additive category equipped with

$$\Sigma : \mathcal{T} \xrightarrow{\sim} \mathcal{T},$$

along with a class of distinguished (exact) triangles,

$$X \xrightarrow{u} Y \xrightarrow{v} Z \xrightarrow{w} \Sigma X \quad (\Delta)$$

satisfying four axioms.

- A **triangulated functor** $F: (\mathcal{T}, \Sigma) \rightarrow (\mathcal{T}', \Sigma')$ satisfies

$$F\Sigma \cong \Sigma'F$$

and $F(\mathcal{T}\text{-exact } \Delta) = \mathcal{T}'\text{-exact } \Delta$.

- Prototypical example:

\mathcal{A} additive category

e.g. $\mathcal{A} = \text{Ab}$
 $\mathcal{A} = R\text{-mod}$
 $\mathcal{A} = \text{Sh}(X)$

$\leadsto \text{Ch}(\mathcal{A})$ additive category of chain complexes in \mathcal{A} .

$\leadsto K(\mathcal{A})$ homotopy category of \mathcal{A} with

$$\text{Ob}(K(\mathcal{A})) = \text{Ob}(\text{Ch}(\mathcal{A})).$$

$$\text{Hom}_{K(\mathcal{A})}(X, Y) = \text{Hom}_{\text{Ch}(\mathcal{A})}(X, Y) / \text{homotopy equivalence}.$$

Then $K(\mathcal{A})$ is triangulated with exact triangles those isomorphic to ones of the form

$$X \xrightarrow{f} Y \longrightarrow \text{cone}(f) \longrightarrow \Sigma X.$$

Also $F: \mathcal{A} \rightarrow \mathcal{A}'$ additive

$\leadsto F: K(\mathcal{A}) \rightarrow K(\mathcal{A}')$ triangulated.

- Goal: If \mathcal{A} is abelian, "localise" $K(\mathcal{A})$ to get triangulated category $D(\mathcal{A})$ which is the "right place" to study left/right exact functors on \mathcal{A} :

- $F: \mathcal{A} \rightarrow \mathcal{A}'$ left exact $\leadsto R^i F: D(\mathcal{A}) \rightarrow D(\mathcal{A}')$ with $H^i R^i F(A) = R^i F(A)$ for $A \in \mathcal{A}$

- Complicated homological facts (e.g. expressed by spectral sequences) about F simplify via $R^i F$.

- $D(\text{Sh}(X))$ will also be the home of perverse sheaves on X .

§2 Localisation [Weibel, 10.3]

- \mathcal{C} a category, $S \subseteq \text{Mor}(\mathcal{C})$.

A localisation of \mathcal{C} at S is a category $S^{-1}\mathcal{C}$ with a functor

$$q: \mathcal{C} \rightarrow S^{-1}\mathcal{C}$$

such that: (1) $q(s) \in \text{Isom}(S^{-1}\mathcal{C})$ for all $s \in S$ ($q(s)$ is an iso)

(2) Any $F: \mathcal{C} \rightarrow \mathcal{D}$ such that $F(s) \in \text{Isom}(\mathcal{D})$ for all $s \in S$ admits a factorisation,

$S^{-1}\mathcal{C}$ unique up to equivalence

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{q} & S^{-1}\mathcal{C} \\ & \searrow F & \downarrow \exists! \\ & & \mathcal{D} \end{array}$$

- If $\text{Ob}(\mathcal{C})$ or S is a set, then $S^{-1}\mathcal{C}$ exists.
More generally, it is a delicate set-theoretic issue, which we will avoid.

- Example: Let $R = \text{comm. ring}$ and consider the category \mathcal{C}_R :

$$\begin{array}{c} \curvearrowright R \\ * \rightarrow * \end{array}$$

$$\text{Hom}(*, *) = R$$

Recall that if $S \subseteq R$ is multiplicatively closed, then we can form $S^{-1}R$.

$$\text{Then } S^{-1}\mathcal{C}_R \cong \mathcal{C}_{S^{-1}R} \quad (\text{exercise})$$

- Example: R and S as above. Then S determines a collection of morphisms

$$S = \{f \in \text{Mor}(R\text{-mod}) : S^{-1}(f) \text{ an iso}\}.$$

$$\text{Then } S^{-1}(R\text{-mod}) \cong (S^{-1}R)\text{-mod.}$$

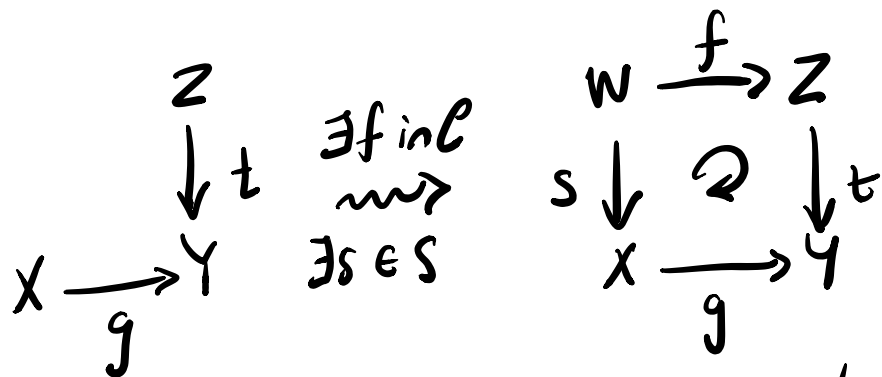
(exercise)

- If S has good properties, we can describe $S^{-1}\mathcal{C}$ explicitly.

- S is called a **multiplicative system** in case:

(1) $1_X \in S$ for all $X \in \mathcal{C}$ and S is closed under composition.

(2) Given a situation with $t \in S$ and $g \in \mathcal{C}$



One condition: " $t^{-1}g = fs^{-1}$ "

moreover, the symmetric statement is true.

(3) Given $X \xrightarrow[g]{f} Y$ in \mathcal{C} ,

$$(\exists t \in S \text{ with } ft = gt) \Leftrightarrow (\exists s \in S \text{ with } sf = sg)$$

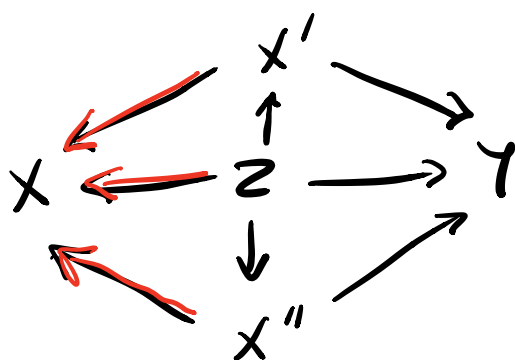
- Theorem [Gabriel-Zisman]: Let S be a (locally small) multiplicative system in \mathcal{C} . Then $S^{-1}\mathcal{C}$ exists and admits an explicit description as follows:

(1) $\text{Ob}(S^{-1}\mathcal{C}) = \text{Ob}(\mathcal{C})$.

(2) $\text{Hom}_{S^{-1}\mathcal{C}}(X, Y) = \left\{ \begin{array}{c} \text{left fractions} \\ X \xleftarrow{s} X' \xrightarrow{f} Y \\ \text{in } S \end{array} \right\} / \sim$

where $(X \leftarrow X' \rightarrow Y) \sim (X \leftarrow X'' \rightarrow Y)$

iff there is a fraction $X \xleftarrow{z} Z \rightarrow Y$ such that the following commutes



Commutative diagram.

- (3) Composition relies on the Ore condition.
(exercise)

- Key situation:

\mathcal{T} = triangulated category
 \mathcal{A} = abelian category.

Then $H: \mathcal{T} \rightarrow \mathcal{A}$ is **cohomological** if it sends exact triangles (Δ) to LES

$$\begin{aligned} \dots &\rightarrow H(\Sigma^i X) \rightarrow H(\Sigma^i Y) \rightarrow H(\Sigma^i Z) \\ &\rightarrow H(\Sigma^{i+1} X) \rightarrow \dots \end{aligned}$$

Notation: $H^i = H \Sigma^i$.

Examples: (1) $H = H^0: K(\mathcal{A}) \rightarrow \mathcal{A}$

(2) $H = \text{Hom}_{\mathcal{T}}(X, -): \mathcal{T} \rightarrow \text{Ab}$ if $X \in \mathcal{T}$.

- Proposition: Let $H: \mathcal{T} \rightarrow \mathcal{A}$ cohomological. If

$S =$ collection of morphisms in \mathcal{T} such that $H^i(s)$ is an iso. for all i ,

Then

(1) S is a multiplicative system

(2) $S^{-1}\mathcal{T}$ is a triangulated category, with

$\mathcal{T} \rightarrow S^{-1}\mathcal{T}$ a triangulated functor.

- Now the derived category $D(\mathcal{A}) = \text{localisation}$ of $K(\mathcal{A})$ arising from $H^0: K(\mathcal{A}) \rightarrow \mathcal{A}$.
(we invert quasi-isomorphisms)

- Alternative construction: $D(\mathcal{A})$ is a "Verdier quotient" of $K(\mathcal{A})$ by the full subcat. generated by acyclic complexes X (i.e. $H^i(X) = 0$)
for all i

[Chambert-Loir]

- Subcategories:

- Bounded below categories $Ch^+(\mathcal{A}), K^+(\mathcal{A}), D^+(\mathcal{A})$
with $H^i(A) = 0$ for $i \ll 0$ for all A .

- Bounded above categories $Ch^-(\mathcal{A}), K^-(\mathcal{A}), D^-(\mathcal{A})$
with $H^i(A) = 0$ for $i \gg 0$ for all A .

- Intersections $Ch^b(\mathcal{A}), K^b(\mathcal{A}), D^b(\mathcal{A})$ of the preceding, the bounded categories.

§3 Derived functors

- \mathcal{A}, \mathcal{B} abelian categories, $F: \mathcal{A} \rightarrow \mathcal{B}$ additive.

$$\rightsquigarrow F: K(\mathcal{A}) \rightarrow K(\mathcal{B}).$$

"uninteresting"

If furthermore F is exact, then

$$\rightsquigarrow F: D(\mathcal{A}) \rightarrow D(\mathcal{B}).$$

- More interesting: what can we say if F is exact on only one side? To fix notation, assume F is left exact.

- Definition: A right derived functor of F is a functor

$$RF: D^+(\mathcal{A}) \rightarrow D^+(\mathcal{B})$$

along with natural transformation η as follows

$$\begin{array}{ccc} K^+(\mathcal{A}) & \xrightarrow{F} & K^+(\mathcal{B}) \\ \eta \downarrow & \swarrow \eta & \downarrow \eta \\ D^+(\mathcal{A}) & \xrightarrow{RF} & D^+(\mathcal{B}) \end{array}$$

such that RF with η is universal in this setup.

- Theorem: If \mathcal{A} has enough injectives, then RF exists.

Sketch proof: Let $\mathcal{I} \subseteq \mathcal{A}$ be full subcategory generated by injective objects. Every $X \in K^+(\mathcal{A})$ admits a quasi-isomorphism

$$X \longrightarrow I \in K^+(\mathcal{I})$$

and also $S^{-1}K^+(\mathcal{I}) \subseteq D^+(\mathcal{A})$ is a fully faithful embedding. This proves that

$$S^{-1}K^+(\mathcal{I}) \cong D^+(\mathcal{A})$$

Since quasi-isomorphisms in $K^+(\mathcal{I})$ are already isomorphisms, $K^+(\mathcal{I}) \cong S^{-1}K^+(\mathcal{I})$, so

$$\Psi: K^+(\mathcal{I}) \cong D^+(\mathcal{A}).$$

Now we take $RF = D^+(\mathcal{A}) \xrightarrow{\Psi^{-1}} K^+(\mathcal{I})$
 $\xrightarrow{F} K^+(\mathcal{B}) \xrightarrow{q} D^+(\mathcal{B})$

So $RF = qF\Psi^{-1}$. Exercise: construct η .

- Theorem: If \mathcal{A} has enough injectives, then

$$H^i R^i F(X) \cong R^i F(X). \quad \leftarrow \text{hypercohomology}$$

In particular, if $X \in \mathcal{A}$, then

$$H^i R^i F(X) \cong R^i F(X).$$

- Examples: $A =$ commutative ring, $\mathcal{A} = A\text{-mod.}$

(1) If M is an A -module, then can form

$$F = \text{Hom}_A(M, -)$$

left exact and

$$H^i R^i F(N) = R^i F(N) = \text{Ext}_A^i(M, N).$$

(2) The functor

$$L = M \otimes_A - : \mathcal{A} \rightarrow \mathcal{A}$$

right exact \rightsquigarrow $LL: D^-(\mathcal{A}) \rightarrow D^-(\mathcal{A})$

is written $LL(N) = M \otimes_A^L N$ and

$$H^{-i} LL(N) = \text{Tor}_i^A(M, N).$$

(Works because \mathcal{A} has enough projectives.)

- Remark: We have $R\text{Hom}(-, -)$ and $- \otimes^L -$ where one input is a module and the other is a complex. A theory of **derived bifunctors** permits both variables to be complexes.